

FEMaLeCoP: OCaml leanCoP and Fairly Efficient Learning

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Talk Overview

- Connection Tableaux and leanCoP
- leanCoP in OCaml and HOL
 - ARLT and Hammers
 - Proof Certification
 - Reconstruction
 - Integration in HOL
- leanCoP and Learning
 - MaLeCoP
 - Features
 - Indexing and Learning
 - Advising

leanCoP: Lean Connection Prover (Jens Otten)

- Connected tableaux calculus
 - **Goal oriented**, good for large theories
- Regularly beats Metis and Prover9 in CASC
 - despite their much larger implementation
 - very good **performance** on some ITP challenges
- **Compact** Prolog implementation, easy to modify
- Has variants for other foundations
 - iLeanCoP

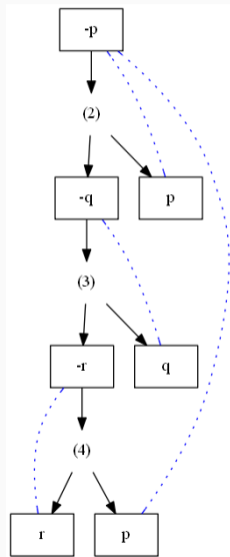
Lean connection Tableaux

	$\overline{\{\}, M, Path}$	<i>Axiom</i>	
	$\frac{C, M, \{\}}{M}$	<i>Start</i>	where $C \in M$, C is positive
	$\frac{C, M, Path \cup \{L_2\}}{C \cup \{L_1\}, M, Path \cup \{L_2\}}$	<i>Reduction</i>	where $\sigma(L_1) = \sigma(\overline{L_2})$
$\frac{C_2 \setminus \{L_2\}, M, Path \cup \{L_1\}}{C \cup \{L_1\}, M, Path}$	$C, M, Path$	<i>Extension</i>	where $\sigma(L_1) = \sigma(\overline{L_2})$, σ is rigid, where $C_1 \in M$, $L_2 \in C_2$, C_2 is a copy of C_1 with vars renamed

leanCoP: Example of Connection Tableau

```
fof(1,conjecture,~p).  
fof(2,axiom,p => q).  
fof(3,axiom,q => r).  
fof(4,axiom,r => ~p).
```

- DNF vs CNF approach (leanCoP vs most resolution provers)
- Axioms \implies Conjecture
- \neg Axioms \vee Conjecture
- CNF for Axioms: (2) $(\neg p \mid q)$ & (3) $(\neg q \mid r)$ & (4) $(\neg r \mid \neg p)$
- DNF for \neg Axioms: (2) $(p \ \& \ \neg q)$ | (3) $(q \ \& \ \neg r)$ | (4) $(r \ \& \ p)$
- starting DNF: (1) $[\neg p]$ (2) $[p, \neg q]$ (3) $[q, \neg r]$ (4) $[r, p]$



leanCoP: Basic Code

```
1 prove ([Lit|Cla], Path, PathLim, Lem, Set) :-
2   %
3   (-NegLit=Lit; -Lit=NegLit) ->
4     ( %
5       %
6         %
7         member (NegL, Path), unify_with_occurs_check (NegL, NegLit)
8       ;
9       lit (NegLit, NegL, Cla1, Grnd1),
10      unify_with_occurs_check (NegL, NegLit),
11      %
12      %
13      %
14      prove (Cla1, [Lit|Path], PathLim, Lem, Set)
15    ),
16    %
17    prove (Cla, Path, PathLim, Lem, Set) .
18 prove ([], _, _, _, _).
```

leanCoP: Actual Code (Optimizations, No history)

```
1 prove ([Lit|Cla], Path, PathLim, Lem, Set) :-
2   \+ (member(LitC, [Lit|Cla]), member(LitP, Path), LitC==LitP),
3   (-NegLit=Lit; -Lit=NegLit) ->
4     (
5       member(LitL, Lem), Lit==LitL
6     ;
7       member(NegL, Path), unify_with_occurs_check(NegL, NegLit)
8     ;
9       lit(NegLit, NegL, Cla1, Grnd1),
10      unify_with_occurs_check(NegL, NegLit),
11      ( Grnd1=g -> true ;
12        length(Path, K), K < PathLim -> true ;
13        \+ pathlim -> assert(pathlim), fail ),
14      prove(Cla1, [Lit|Path], PathLim, Lem, Set)
15    ),
16    ( member(cut, Set) -> ! ; true ),
17    prove(Cla, Path, PathLim, [Lit|Lem], Set) .
18 prove([], _, _, _, _).
```

Automated Reasoning in Large Theories

- Prove goals **automatically** in large formal theories
 - ATP translation of MML (\approx 50k proofs today)
 - Isabelle/HOL (\approx 60k proofs today)
 - HOL Light/Flyspeck (\approx 30k proofs today)
 - More proof assistant corpora: HOL4, ACL2, Coq
- **Useful** for ITP
 - Sledgehammer, HOL(y)Hammer, MizAR
 - But needs Reconstruction

[Urban03,...]

[Paulson05,Blanchette]

[KU12]

Request Advice:

Input the HOL Light formula to prove and select HOL Light session:

-
-

(cache:OK)(session:OK)(parse:OK)SSSSAWAAWAW

Result (3.81s): CONVEX_RELATIVE_INTERIOR POLYHEDRON_IMP_CONVEX

Replaying: SUCCESS (0.29s):SIMP_TAC[POLYHEDRON_IMP_CONVEX;CONVEX_RELATIVE_INTERIOR]

Existing Proof Reconstruction

- **General** ATP search tactics producing ITP proof objects:
 - Metis (Isabelle, HOL4)
 - MESON, Prover9 (HOL Light)
 - Mizar by
- Parse **TSTP/SMT** proofs
 - Create subgoals that match ATP intermediate steps
 - Automatically solve all simple subgoals
 - **Skolemization** of type variables is an issue
- The smarter ATPs we can integrate in ITPs, the better
 - Not just for the Hammers
- Need for speed
 - Thousands of reconstruction steps in ITP projects
- ATP proof search blowup

Rewriting leanCoP in OCaml

- Parts of the Prolog technology missing in functional languages
- Use an **explicit stack** for keeping track of the current proof state (including the trail of variable bindings)
 - In the main `prove` function we add explicit arguments:
`stack (stack)`, `subst (trail)` and `off (offset in the trail)`
 - The stack keeps a list of tuples that are given as arguments to the recursive invocations of `prove`
 - When a proof is found, the exception 'Solved' is raised and the function exits with this exception.
- An alternative would be to use the **continuation passing style**

OCaml code (No history)

```
1 let rec prove path lim lem stack = function (lit :: cla) ->
2   if not (exists2 eq path (lit :: cla)) then
3     let neglit = negate lit in
4     if not (exists (substeq lit) lem && (prove path lim lem stack cla; cut)) then
5     if not (fold_left (fun sf plit -> sf ||
6       try (unify_lit neglit plit; prove path lim (lit :: lem) stack cla; cut)
7       with Unify -> sf) false path) then
8     let iter_fun (lit2, cla2, ground) =
9       if lim > 0 || ground then
10        try let cla1 = unify_rename (snd lit) (lit2, cla2) in
11          prove (lit :: path) (lim - 1) lem ((if cut then lim else -1),
12            path, lim, lit :: lem, cla) :: stack) cla1 with Unify -> () in
13        try iter iter_fun (try assoc neglit lits with Not_found -> [])
14        with Cut n -> if n = lim then () else raise Cut n))
15 | [] -> match stack with
16   (ct, path, lim, lem, cla) :: t ->
17     prove path lim lem t cla; if ct > 0 raise (Cut ct)
18 | [] -> raise Solved;;
```

Differences

- Implementation of **Prolog cut** (!) in OCaml:
 - different mechanism in each of the three cases
- The Prolog code is **elegant**
 - But the OCaml code is a bit more efficient.
- A simple `List.exists` call is enough for finding a lemma
 - no need to backtrack
- For equality checking and unification under substitution we reuse **MESON code**:
 - substitutions as association lists
 - applications of substitutions are delayed until an equality check or a unification step

Eval I: HOL Light MESON calls without splitting (872 goals, 5s) ¹

Prover	Theorem (%)	Unique
OcaML-leanCoP (cut)	759 (87.04)	2
OcaML-leanCoP (nocut)	759 (87.04)	2
Prolog-leanCoP (cut)	752 (86.23)	0
Prolog-leanCoP (nocut)	751 (86.12)	0
Metis (2.3)	708 (81.19)	26
Meson	683 (78.32)	4
any	832 (95.41)	

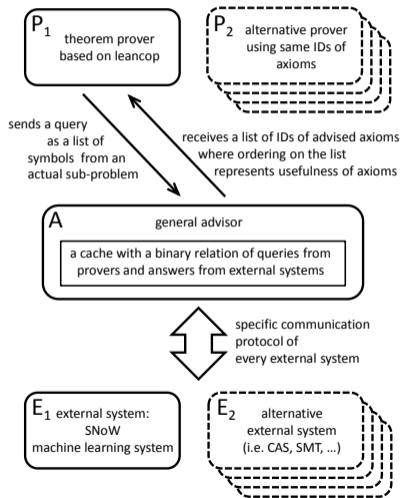
¹Evaluation outside HOL Light

Reconstruction of leanCoP proofs in HOL Light

- Transformation from **HOL** to **FOL** and clausification
 - Tactics reuse Harrison's MESON code
 - Needs to preserve leanCoP's goal-directed approach
 - The conjecture is separated from the axioms
- All transformations are done on the **CNF** rather than **DNF**
 - The two are dual
- MESON's **proof reconstruction** needs modification for the use of **lemmas**

Very secure HOL Light certification of leanCoP proofs

Learning: MaLeCoP



FEMaLeCoP: Advice Overview and Used Features

- Advise the:
 - selection of clause for every tableau extension step
- Proof state: weighted vector of **symbols** (or terms)
 - extracted from all the literals on the active path
 - Frequency-based weighting (IDF)
 - Simple **decay factor** (using maximum)
- **Consistent clausification**
 - formula $?[X] : p(X)$ becomes $p('skolem(?[A] : p(A), 1)')$
- Advice using custom sparse **naive Bayes**
 - association of the features of the proof states
 - with contrapositives used for the successful extension steps

FEMaLeCoP: Data Collection and Indexing

- Slight extension of the saved proofs
 - Training Data: pairs (path, used extension step)
- **External** Data Indexing (incremental)
 - `te_num`: number of training examples
 - `pf_no`: hashtable from features to number of occurrences $\in \mathbb{Q}$
 - `cn_no`: hashtable from contrapositives to numbers of occurrences
 - `cn_pf_no`: hashtable of maps of `cn/pf` co-occurrences
- **Problem Specific** Data
 - Upon start FEMaLeCoP reads
 - **only current-problem** relevant parts of the training data
 - `cn_no` and `cn_pf_no` filtered by contrapositives in lit matrix
 - `pf_no` and `cn_pf_no` filtered by possible features in the problem

Naive Bayes

If more than one possible extension step

Estimate relevance of each contrapositive cn by

$$\sigma_1 \ln t + \sum_{f \in (\bar{f} \cap \bar{s})} i(f) \ln \frac{\sigma_2 s(f)}{t} + \sigma_3 \sum_{f \in (\bar{f} - \bar{s})} i(f) + \sigma_4 \sum_{f \in (\bar{s} - \bar{f})} i(f) \ln \left(1 - \frac{s(f)}{t}\right)$$

where

- \bar{f} are the features of the path
- \bar{s} are the features that co-occurred with cn
- $t = cn_no(cn)$
- $s = cn_fp_no(cn)$
- i is the IDF
- σ_* are experimentally chosen parameters

It cannot work?

Inference speed ...

It cannot work? But it does!

Inference speed ... drops to about 40%, but:

Prover	Proved (%)
OCaml-leanCoP	574 (27.6%)
FEMaLeCoP	635 (30.6%)
together	664 (32.0%)

(MPTP bushy problems, 60 s)

Summary

- OCaml version of leanCoP
 - outperforms Metis and MESON, sometimes very significantly
- Reconstruction of leanCoP proofs in HOL Light
 - Useful as reconstruction component of HOL(y)Hammer
 - Certification of leanCoP TPTP proofs
- Learning
 - Three levels of indexing
 - Proper integration
- Future Work:
 - Strategies, more evaluation, more HOL-based ITPs
 - Intuitionistic version
 - More learning algorithms