MACHINE LEARNING AND AUTOMATED REASONING - INTRODUCTION

Josef Urban

Czech Technical University in Prague

March 8, 2019





European Research Council Established by the European Commission

Course Overview

- Connections between two AI fields: Machine Learning (ML) and Automated Reasoning (AR)
- ML: apply various forms of *inductive reasoning* to large datasets to obtain the most plausible explanations, models and conjectures
- AR: apply various forms of *deductive reasoning* to prove that particular explanations and conjectures are correct.
- · Humans combine induction and deduction let's teach computers too!
- · We will mostly explore ML/AR combinations in a formal proof setting
- Typical problem: How can learning help with logical reasoning?

Course Overview - Particular settings and topics

- ML and first-order logic (FOL), saturation-style theorem provers (ATPs)
- Higher-order logic (HOL), Set theory, formal proof assistants (ITPs)
- ML and reasoning in large theories, hammers for ITP, premise selection
- · Symbolic vs statistical learning for theorem proving
- ML in tableau-style and tactical reasoning systems
- Learning in propositional logic (SAT), QBF, SMT, instantiation-based methods and model finding.
- Representations and conjecturing how do we characterize reasoning data for learning?
- Feedback loops for proving and learning, reinforcement learning of ATP, positive/negative proof mining
- Alignment and translation between informal and formal corpora, automated formalization
- Exam: do a small project in combining ML and AR

Induction/Learning vs Reasoning - Henri Poincaré



- Science and Method: Ideas about the interplay between correct deduction and induction/intuition
- "And in demonstration itself logic is not all. The true mathematical reasoning is a real induction [...]"
- I believe he was right: strong general reasoning engines have to combine deduction and induction (learning patterns from data, making conjectures, etc.)

Learning vs Reasoning – Alan Turing 1950 – Al



- 1950: Computing machinery and intelligence AI, Turing test
- "We may hope that machines will eventually compete with men in all purely intellectual fields." (regardless of his 1936 undecidability result!)
- last section on Learning Machines(!):
- "But which are the best ones [fields] to start [learning on] with?"
- "... Even this is a difficult decision. Many people think that a very abstract activity, like the playing of chess, would be best."
- · Why not try with large computer-understandable math corpora?

Intuition vs Formal Reasoning - Poincaré vs Hilbert



[Adapted from: Logicomix: An Epic Search for Truth by A. Doxiadis]

What is Formal Mathematics?

- · Developed thanks to the Leibniz/Russell/Frege/Hilbert/... program
- Mathematics put on formal logic foundations (symbolic computation)
- ... which btw. led also to the rise of computers (Turing/Church, 1930s)
- Formal math (1950/60s): combine formal foundations and the newly available computers
- Conceptually very simple:
- · Write all your axioms and theorems so that computer understands them
- · Write all your inference rules so that computer understands them
- · Use the computer to check that your proofs follow the rules
- But in practice, it turns out not to be so simple
- · Many approaches, still not mainstream, but big breakthroughs recently

tiny proof from Hardy & Wright:

Theorem 43 (Pythagoras' theorem). $\sqrt{2}$ is irrational. The traditional proof ascribed to Pythagoras runs as follows. If $\sqrt{2}$ is rational, then the equation

$$a^2 = 2b^2$$
 (4.3.1)

is soluble in integers *a*, *b* with (a, b) = 1. Hence a^2 is even, and therefore *a* is even. If a = 2c, then $4c^2 = 2b^2$, $2c^2 = b^2$, and *b* is also even, contrary to the hypothesis that (a, b) = 1.

Irrationality of $\sqrt{2}$ (Formal Proof Sketch)

exactly the same text in Mizar syntax:

```
theorem Th43: :: Pythagoras' theorem
  sqrt 2 is irrational
proof
  assume sqrt 2 is rational;
  consider a,b such that
4 3 1: a^2 = 2 \cdot b^2 and
    a,b are relative prime;
  a^2 is even;
  a is even;
  consider c such that a = 2 * c;
  4 \star c^2 = 2 \star b^2;
  2 \star c^2 = b^2;
  b is even;
  thus contradiction;
end;
```

Irrationality of $\sqrt{2}$ in HOL Light

let SQRT_2_IRRATIONAL = prove (`~rational(sqrt(&2))`, SIMP_TAC[rational; real_abs; SQRT_POS_LE; REAL_POS] THEN REWRITE_TAC[NOT_EXISTS_THM] THEN REPEAT GEN_TAC THEN DISCH_THEN(CONJUNCTS_THEN2 ASSUME_TAC MP_TAC) THEN SUBGOAL_THEN `~((&p / &q) pow 2 = sqrt(&2) pow 2)` (fun th -> MESON_TAC[th]) THEN SIMP_TAC[SQRT_POW_2; REAL_POS; REAL_POW_DIV] THEN ASM_SIMP_TAC[REAL_EQ_LDIV_EQ; REAL_OF_NUM_LI; REAL_POW_LT; ARITH_RULE `0 < q <=> ~(q = 0)`] THEN ASM_MESON_TAC[NSQRT_2; REAL_OF_NUM_POW; REAL_OF_NUM_MUL; REAL_OF_NUM_EQ]);;

Irrationality of $\sqrt{2}$ in Isabelle/HOL

```
theorem sgrt2 not rational:
  "sort (real 2) ∉ 0"
proof
 assume "sqrt (real 2) \in \mathbb{Q}"
  then obtain m n :: nat where
    n_nonzero: "n \neq 0" and sqrt_rat: "!sqrt (real 2)! = real m / real n"
    and lowest_terms: "gcd m n = 1" ...
 from n_nonzero and sqrt_rat have "real m = {sqrt (real 2)} * real n" by simp
  then have "real (m^2) = (sort (real 2))^2 * real <math>(n^2)"
    by (auto simp add: power2 eg square)
  also have "(sgrt (real 2))<sup>2</sup> = real 2" by simp
  also have "... * real (m^2) = real (2 * n^2)" by simp
  finally have eq: m^2 = 2 * n^2 ...
  hence "2 dvd m<sup>2</sup>"...
  with two is prime have dvd m: "2 dvd m" by (rule prime dvd power two)
  then obtain k where "m = 2^* k"
  with eq have "2 * n^2 = 2^2 * k^2" by (auto simp add: power2 eq square mult ac)
  hence "n^2 = 2 * k^2" by simp
  hence "2 dvd n^2"...
  with two_is_prime have "2 dvd n" by (rule prime_dvd_power_two)
  with dvd m have "2 dvd qcd m n" by (rule qcd greatest)
  with lowest terms have "2 dvd 1" by simp
 thus False by arith
ged
```

Irrationality of $\sqrt{2}$ in Coq

```
Theorem irrational_sqrt_2: irrational (sqrt 2%nat).
intros p q H H0; case H.
apply (main_thm (Zabs_nat p)).
replace (Div2.double (q * q)) with (2 * (q * q));
[idtac | unfold Div2.double; ring].
case (eq_nat_dec (Zabs_nat p * Zabs_nat p) (2 * (q * q))); auto; intros H1.
case (not_nm_INR _ _ H1); (repeat rewrite mult_INR).
rewrite <- (sqrt_def (INR 2)); auto with real.
rewrite H0; auto with real.
assert (q <> 0%R :> R); auto with real.
field; auto with real; case p; simpl; intros; ring.
Oed.
```

Irrationality of $\sqrt{2}$ in Metamath

\${

```
$d x y $.
$( The square root of 2 is irrational. $)
sqr2irr $p |- ( sqr ` 2 ) e/ QQ $=
```

(vx vy c2 csqr cfv cq wnel wcel wn cv cdiv co wceq cn wrex cz cexp cmulc sqr2irrlem3 sqr2irrlem5 bi2rexa mtbir cc0 clt wbr wa wi wb nngtOt adantr cr axOre ltmuldivt mp3an1 nnret zret syl2an mpd ancoms 2re 2pos sqrgtOi breq2 mpbii syl5bir cc nncnt mulzer2t syl breq1d adantl sylibd exp r19.23adv anc21i elnnz syl6ibr impac r19.22i2 mto elq df-nel mpbir) CDEZFGWDFHZIWEWDAJZBJZKLZMZBNOZAPOZWKWJANOZWLWFCQLCWGCQLRLMZBNOANOABSWIWM ABNNWFWGTUAUBWJWJAPNWFPHZWJWFNHZWNWJNNUCWFUDUEZUFWOWNWJWPNWNIWPBNWNWGMHZW IWPUGWNWQUFZWIUCGRLZWFUDUEZWPWRWTUCWHUDUEZWIWQWNWTXAUHZWQWNUFUCWGUDUEZXB WQXCWNWGUIUJWGUKHZWFUKHZXCXBUGZWQWNUCUKHXDXEXFULUCWGWFUMUNWGUOWFUPUQURUSW IUCWDUDUEXACUTVAVBWDHHUCUDVCVDVEWQWTWPUHWNWQWSUCWFUDUWQWGVFHWSUCMWGVGWGVHV IVJVKVLVMVNOWFVPVQVRVSVTABWDWAUBWDFWBWC \$.

\$([8-Jan-02] \$)

\$}

Irrationality of $\sqrt{2}$ in Metamath Proof Explorer

ဓ 🐵 sqr2irr - Metamath Proof Explorer - Chromium

💦 sqr2irr - Metamat × 👅

< > 😋 🗈 us.metamath.org/mpegif/sqr2irr.html

0.5		AIP	Ξ
~ (~)	20,	~	_

			Proof of Theorem sqr2irr
Step	Нур	Ref	Expression
1		sqr2irrlem3 10838	$\dots \land \vdash \neg \exists x \in \mathbb{N} \exists y \in \mathbb{N} (x \uparrow 2) = (2 \cdot (y \uparrow 2))$
2		sqr2irrlem5 10840	$\dots \leftarrow \vdash ((x \in \mathbb{N} \land y \in \mathbb{N}) \rightarrow ((\sqrt{2}) = (x / y) \leftrightarrow (x^{\uparrow}2) = (2 \cdot (y^{\uparrow}2))))$
3	2	2rexbiia 2329	$\dots : \vdash (\exists x \in \mathbb{N} \exists y \in \mathbb{N} (\sqrt{2}) = (x / y) \leftrightarrow \exists x \in \mathbb{N} \exists y \in \mathbb{N} (x \uparrow 2) = (2 \cdot (y \uparrow 2)))$
4	1, 3	mtbir 288	$\dots + \vdash \neg \exists \mathbf{x} \in \mathbb{N} \exists \mathbf{y} \in \mathbb{N} (\sqrt{2}) = (\mathbf{x} / \mathbf{y})$
5		2 <u>re</u> 8838	
6		2pos 8849	12 H 0 < 2
7	<u>5, 6</u>	sqrgt0ii 10213	·····································
8		breq2 3595	$\dots \dots \mapsto \vdash ((\sqrt{2}) = (x / y) \rightarrow (0 < (\sqrt{2}) \leftrightarrow 0 < (x / y)))$
9	<u>7, 8</u>	mpbii 200	$\dots \dots \dots \mapsto \vdash ((\sqrt{2}) = (x / y) \to 0 < (x / y))$
10		ZTC 9029	$\dots \dots \dots \square 12 \vdash (x \in \mathbb{Z} \to x \in \mathbb{R})$
11	10	adantr 444	$\dots \dots \mapsto \vdash ((x \in \mathbb{Z} \land y \in \mathbb{N}) \rightarrow x \in \mathbb{R})$
12		nnre \$788	$\dots \dots \square \vdash (y \in \mathbb{N} \rightarrow y \in \mathbb{R})$
13	12	adantl 445	$\dots \dots \mapsto \vdash ((x \in \mathbb{Z} \land y \in \mathbb{N}) \rightarrow y \in \mathbb{R})$
14		nngt0 8807	$\dots \dots \square \vdash (\mathbf{y} \in \mathbb{N} \rightarrow 0 < \mathbf{y})$
15	14	adantl 445	$\dots \dots \mapsto \vdash ((x \in \mathbb{Z} \land y \in \mathbb{N}) \rightarrow 0 < y)$
16		gt0div sess	$\dots \dots \mapsto \vdash ((x \in \mathbb{R} \land y \in \mathbb{R} \land 0 < y) \rightarrow (0 < x \leftrightarrow 0 < (x / y)))$
17	11, 13, 15, 16	syl3anc 1145	10 \vdash $((x \in \mathbb{Z} \land y \in \mathbb{N}) \rightarrow (0 < x \leftrightarrow 0 < (x / y)))$
18	<u>9, 17</u>	syl5ibr 210	$\dots \dots \oplus \vdash ((x \in \mathbb{Z} \land y \in \mathbb{N}) \to ((\sqrt{2}) = (x / y) \to 0 < x))$
19		simpl 436	$\dots \dots \dots \dots \dots \mapsto \vdash ((x \in \mathbb{Z} \land y \in \mathbb{N}) \to x \in \mathbb{Z})$
20	<u>18, 19</u>	jctild 522	$\dots \dots \Vdash \vdash ((x \in \mathbb{Z} \land y \in \mathbb{N}) \to ((\sqrt{2}) = (x / y) \to (x \in \mathbb{Z} \land 0 < x)))$
21		elnnz 9035	$\dots \dots \Vdash (x \in \mathbb{N} \leftrightarrow (x \in \mathbb{Z} \land 0 < x))$
22	<u>20, 21</u>	<u>syl6ibr</u> 216	$\dots, \forall \vdash ((x \in \mathbb{Z} \land y \in \mathbb{N}) \rightarrow ((\sqrt{2}) = (x / y) \rightarrow x \in \mathbb{N}))$
23	22	rexlimdva 2414	$\dots \land \land \vdash (x \in \mathbb{Z} \rightarrow (\exists y \in \mathbb{N} (\sqrt{2}) = (x / y) \rightarrow x \in \mathbb{N}))$
24	23	impac 598	$\ldots s \vdash ((x \in \mathbb{Z} \land \exists y \in \mathbb{N} (\sqrt{2}) = (x / y)) \rightarrow (x \in \mathbb{N} \land \exists y \in \mathbb{N} (\sqrt{2}) = (x / y)))$
25	24	reximi2 2396	$\dots \models \vdash (\exists x \in \mathbb{Z} \exists y \in \mathbb{N} (\sqrt{2}) = (x / y) \rightarrow \exists x \in \mathbb{N} \exists y \in \mathbb{N} (\sqrt{2}) = (x / y))$
26	4, <u>25</u>	<u>mto</u> 165	$ \vdash \neg \exists x \in \mathbb{Z} \exists y \in \mathbb{N} (\sqrt{2}) = (x / y)$
27		elq 9308	$\square \exists \vdash ((\checkmark' 2) \in \mathbb{Q} \leftrightarrow \exists x \in \mathbb{Z} \exists y \in \mathbb{N} (\checkmark' 2) = (x / y))$
28	26, 27	mtbir 288	$z \vdash \neg (\sqrt{2}) \in \mathbb{Q}$
29		df-nel 2210	$2 \vdash ((\sqrt{2}) \notin \mathbb{Q} \leftrightarrow \neg (\sqrt{2}) \in \mathbb{Q})$
30	28, 29	mpbir 198	i⊢ (√'2) ∉ Q

Colors of variables: wff set class

Today: Computers Checking Large Math Proofs



Big Example: The Flyspeck project

 Kepler conjecture (1611): The most compact way of stacking balls of the same size in space is a pyramid.



- Formal proof finished in 2014
- · 20000 lemmas in geometry, analysis, graph theory
- All of it at https://code.google.com/p/flyspeck/
- · All of it computer-understandable and verified in HOL Light:
- polyhedron s /\ c face_of s ==> polyhedron c
- However, this took 20 30 person-years!

top 100 of interesting theorems/proofs (Paul & Jack Abad, 1999, tracked by Freek Wiedijk)

- 1. √2 ∉ ℚ
- 2. fundamental theorem of algebra
- 3. $|\mathbb{Q}| = \aleph_0$

4.
$$a \bigsqcup_{b}^{c} \Rightarrow a^{2} + b^{2} = c^{2}$$

5.
$$\pi(x) \sim \frac{x}{\ln x}$$

- 6. Gödel's incompleteness theorem
- 7. $\left(\frac{p}{q}\right)\left(\frac{q}{p}\right) = (-1)^{\frac{p-1}{2}\frac{q-1}{2}}$
- 8. impossibility of trisecting the angle and doubling the cube
- 32. four color theorem
- 33. Fermat's last theorem
- 99. Buffon needle problem
- 100. Descartes rule of signs

top 100 of interesting theorems/proofs (Paul & Jack Abad, 1999, tracked by Freek Wiedijk)

- 1. √2 ∉ ℚ
- 2. fundamental theorem of algebra
- 3. $|\mathbb{Q}| = \aleph_0$

4.
$$a = b^{c} \Rightarrow a^{2} + b^{2} = c^{2}$$

- 5. $\pi(x) \sim \frac{x}{\ln x}$
- 6. Gödel's incompleteness theorem
- 7. $\left(\frac{p}{q}\right)\left(\frac{q}{p}\right) = (-1)^{\frac{p-1}{2}\frac{q-1}{2}}$
- 8. impossibility of trisecting the angle and doubling the cube
- 32. four color theorem
- 33. Fermat's last theorem
- 99. Buffon needle problem
- 100. Descartes rule of signs

all together	88%
HOL Light	86%
Mizar Isabelle Coq ProofPower	57% 52% 49% 42%
Metamath	24%

- ACL2 18%
 - PVS 16%

top 100 of interesting theorems/proofs (Paul & Jack Abad, 1999, tracked by Freek Wiedijk)

1. $\sqrt{2} \notin \mathbb{Q}$	allt
fundamental theorem of algebra	Ц
3. $ \mathbb{Q} = \aleph_0$	
4. $a b^{c} \Rightarrow a^{2} + b^{2} = c^{2}$	
5. $\pi(x) \sim \frac{x}{\ln x}$	
6. Gödel's incompleteness theorem	_
7. $\left(\frac{p}{q}\right)\left(\frac{q}{p}\right) = (-1)^{\frac{p-1}{2}\frac{q-1}{2}}$	Pro
8. impossibility of trisecting the	Me
angle and doubling the cube	
32. four color theorem	
33. Fermat's last theorem	
99. Buffon needle problem	
100. Descartes rule of signs	

all together	88%
HOL Light	86%
Mizar	57%
Isabelle	52%
Coq	49%
ProofPower	42%
Metamath	24%
	18%

PVS

16%

top 100 of interesting theorems/proofs (Paul & Jack Abad, 1999, tracked by Freek Wiedijk)

- 1. √2 ∉ ℚ
- 2. fundamental theorem of algebra
- 3. $|\mathbb{Q}| = \aleph_0$

4.
$$a = b^{c} \Rightarrow a^{2} + b^{2} = c^{2}$$

- 5. $\pi(x) \sim \frac{x}{\ln x}$
- 6. Gödel's incompleteness theorem
- 7. $\left(\frac{p}{q}\right)\left(\frac{q}{p}\right) = (-1)^{\frac{p-1}{2}\frac{q-1}{2}}$
- 8. impossibility of trisecting the angle and doubling the cube
- 32. four color theorem
- 33. Fermat's last theorem
- 99. Buffon needle problem
- 100. Descartes rule of signs

- all together 88% HOL Light 86% Mizar 57% Isabelle 52% Coq 49% ProofPower 42% Metamath 24%
 - ACL2 18% PVS 16%

Named Theorems in the Mizar Library

	Be FM - Chromium			
4	> C n.uwb.edu.pl/	- mmlquery/fillin.php?filledfilename=mml-facts.mqt&argument=number+102	ପ୍	☆ 🗔 🐵 ≡
	Mizar home, download files: abstr. articles	The most important facts in MML (<u>decode</u>)	add de	scription
	bin, doc, emacs gabs.	See also Name carrying facts/theorems/definitions in MML		1.00.000
	fmbibs, gabs (more)	1 "Alexander\'s Lemma"	=> <u>WAYBEL_7:31</u>	VOTE
	semantic MML	2 "All Primes (1 mod 4) Equal the Sum of Two Squares"	=> <u>NAT_5:23</u>	VOTE
		3 "Axiom of Choice"	=> WELLORD2:18	VOTE
		4 "Baire Category Theorem (Banach spaces)"	=> <u>LOPBAN_5:3</u>	VOTE
	情報	5 "Baire Category Theorem (Hausdorff spaces)"	=> <u>NORMSP_2:10</u>	VOTE
	TELER	6 "Baire Category Theorem for Continuous Lattices"	=> WAYBEL12:39	VOTE
	MML Query (beta)	7 "Banach Fix Point Theorem for Compact Spaces"	=> <u>ALI2:1</u>	VOTE
	Template maker	8 "Banach-Steinhaus theorem (uniform boundedness)"	=> <u>LOPBAN_5:7</u>	VOTE
	Environment explanation	9 "Bertrand\'s Ballot Theorem"	=> <u>BALLOT_1:28</u>	VOTE
		10 "Bertrand\'s postulate"	=> <u>NAT_4:56</u>	VOTE
	Mizar TWiki	11 "Bezout\'s Theorem"	=> NEWTON:67	VOTE
	Megrez services	12 "Bing Theorem"	=> <u>NAGATA_2:22</u>	VOTE
	Journals:	13 "Binomial Theorem"	=> BINOM:25	VOTE
	FM: MetaPRESS,	14 "Birkhoff Variety Theorem"	=> BIRKHOFF:sch 12	VOTE
	server, proof-read,	15 "Bolzano theorem (intermediate value)"	=> TOPREAL5:8	VOTE
	MM&A	16 "Bolzano-Weierstrass Theorem (1 dimension)"	=> <u>SEO_4:40</u>	VOTE
	(preparation)	17 "Borsuk Theorem on Decomposition of Strong Deformation Retracts"	=> BORSUK 1:42	VOTE
	Constant and I band	18 "Borsuk-Ulam Theorem"	=> BORSUK 7:condreg 3	VOTE
	Downloads	19 "Boundary Points of Locally Euclidean Spaces"	=> MFOLD 0:2	VOTE
		20 "Brouwer Fixed Point Theorem"	=> BROUWER:14	VOTE
	Mizar syntax, xml, txt	21 "Brouwer Fixed Point Theorem for Disks on the Plane"	=> BROUWER:15	VOTE
	MML 5 25 1220	22 "Brouwer Fixed Point Theorem for Intervals"	=> TREAL 1:24	VOTE
	- most important facts	23 "Brown Theorem"	=> GCD 1:40	VOTE
	(other collection)	24 "Cantor Theorem"	=> CARD 1:14	VOTE
	Birkhoff	25 "Cantor-Bernstein Theorem"	=> CARD 1:10	VOTE -

- · Kepler Conjecture (Hales et all, 2014, HOL Light, Isabelle)
- Feit-Thompson (odd-order) theorem
 - Two graduate books
 - · Gonthier et all, 2012, Coq
- Compendium of Continuous Lattices (CCL)
 - · 60% of the book formalized in Mizar
 - · Bancerek, Trybulec et all, 2003
- The Four Color Theorem (Gonthier and Werner, 2005, Coq)

Mid-size Formalizations

- Gödel's First Incompleteness Theorem Natarajan Shankar (NQTHM), Russell O'Connor (Coq)
- Brouwer Fixed Point Theorem Karol Pak (Mizar), John Harrison (HOL Light)
- Jordan Curve Theorem Tom Hales (HOL Light), Artur Kornilowicz et al. (Mizar)
- Prime Number Theorem Jeremy Avigad et al (Isabelle/HOL), John Harrison (HOL Light)
- Gödel's Second incompleteness Theorem Larry Paulson (Isabelle/HOL)
- Central Limit Theorem Jeremy Avigad (Isabelle/HOL)

Large Software Verifications

- seL4 operating system microkernel
 - · Gerwin Klein and his group at NICTA, Isabelle/HOL
- CompCert a formally verified C compiler
 - · Xavier Leroy and his group at INRIA, Coq
- EURO-MILS verified virtualization platform
 - ongoing 6M EUR FP7 project, Isabelle
- CakeML verified implementation of ML
 - Magnus Myreen, HOL4

Central Limit Theorem in Isabelle/HOL



Sylow's Theorems in Mizar

```
theorem :: GROUP_10:12
for G being finite Group, p being prime (natural number)
holds ex P being Subgroup of G st P is_Sylow_p-subgroup_of_prime p;
theorem :: GROUP_10:14
for G being finite Group, p being prime (natural number) holds
  (for H being Subgroup of G st H is_p-group_of_prime p holds
    ex P being Subgroup of G st
    P is_Sylow_p-subgroup_of_prime p & H is Subgroup of P) &
  (for P1.P2 being Subgroup of G
    st P1 is_Sylow_p-subgroup_of_prime p & P2 is_Sylow_p-subgroup_of_prime p
    holds P1.P2 are_conjugated);
```

```
theorem :: GROUP_10:15
```

```
for G being finite Group, p being prime (natural number) holds
  card the_sylow_p-subgroups_of_prime(p,G) mod p = 1 &
  card the_sylow_p-subgroups_of_prime(p,G) divides ord G;
```

Gödel Theorems in Isabelle









Implementing Prover Trident for SL, Stockholm

In this project, Prover Technology provides the Prover Trident solution to Ansaldo STS, for development and safety approval of interlocking software for Roslagsbanan, a mainline railway line that connects...



Formal Verification of SSI Software for NYCT, New York

New York City Transit (NYCT) is modernizing the signaling system in its subway by installing CBTC and replacing relay-based interlockings with computerized, solid state interlockings (SSIs).



Our Formal Verification Solution for RATP, Paris

In this project Prover Technology collaborated with RATP in creating a formal verification solution to meet RATP demand for safety verification of interlocking software. RATP had selected a computerized...

● Applications Places ● 図して/NS Unhackable × \\ REMS	× \vee jt Robots chai × \vee Sp Startpage \vee x \vee Sp byron co	ok × $\sqrt{10}$ Byron Cook × $\sqrt{50}$ AWS	S Securi 🛛 🍎 Automat	Tue 21:15 Tue 21:15 ted ×
← → C Secure https://aws.amazon.com	/blogs/security/tag/automated-reasoning/		☆ [@]	• 🖻 🤉 👷 •
aws			Account 👻 Sign	n Up
Products Solutions Pricing Learn P	artner Network AWS Marketplace Explore More Q			
Blog Home Category - Edition - Follo	w 🗸		Search Blogs	Q
Tag: Automated reasoning	I			
	How AWS SideTrail verifies key AWS cryptography code by Dariel Schwartz-Narbonne on 15 OCT 2018 in Security, Identity, We know you want to spend your time learning valuable ne vorying about managing infrastructure. That's why we're al services, particularly when it comes to cloud security. With th Read More	& Compliance Permaink Comments v skills, building innovative software, a ways looking for ways to help you aut hat in mind, we recently developed [Share and scaling up applications tomate the management o .]	s — not if AWS
Next Gen Cloud Security with Automated Reasoning aWS podcast	Podcast: Al tech named automated reasoning provides n by Suppin Anami (an 08 OCT 2018) in Security, Identity, & Complian AWS just released a new podcast on how next generation is elvels of assurance for key components of your AWS archite discusses how automated reasoning is embedded within AW Read More	ext-gen cloud security a Pemalink Comments Share curity technology, backed by automati- ture. Byron Cook, Director of the AWS S services and code and the tools cust	ed reasoning, is providing) 5 Automated Reasoning Gri 30mers can []	you higher oup,
	Daniel Schwartz-Narbonne shares how automated reaso by Suphya Anand on 02 OCT 2018 in Security, Security, Identity, & O I recently sait down with Daniel Schwartz-Narbonne, a softw WS, to learn more about the groundbreaking work host his technology based on mathematical logic, to prove that key o Read More	ning is helping achieve the provable impliance Permaink ● Comments ● 5 are development engineer in the Auto is doing in cloud security. The team i components of the cloud are operating	e security of AWS boot co hare imated Reasoning Group (A uses automated reasoning, g as []	ode ARG) at a

Applications Places Applications Places NS ×	× \[it] × \[sp	× (sp × (in	× ⟨sp × ∖ e	× (sp × (w ×)	(= × (sp ×) G	× (83 × (83 × (89 ×	(so x (so x	U = 1	ue 21:40 Tu 4 × \	Josef
← → C ≜ Secure	https://www.abs	int.com/compce	rt/					२ 🏠 🥯	8 <mark>0</mark> 8	N 🔒
🗳 Absint		Sup	port	News	About us	Contact	Search			Î
	CompCert	How it works	New in 18.10	Try now						
A variant of record.goto counting the number of L	that alle incre nag instructions	nentally compt i starting at i	Formal	lly verifie	ed compil	ation	ivader († 19.			
Definition measure.edge to	n 13.43 lips 4. mai	ie) tij, maake -o	-att : +add -> +at		Lenna Speket.pr	termi.				
fion a ro of pay Waapy and if pay Waapy and if pay Waapy and find the field of the	CompCert is a written in C ai extensions. It	formally verified formally verified of meeting high produces mach	ed optimizing C co n levels of assuran nine code for ARM,	mpiler. Its intended ce. It accepts most o PowerPC, x86, and	use is compiling safe f the ISO C 99 langua RISC-V architectures.	ty-critical and mission-critic ge, with some exceptions a	al software ind a few			
Befinition record.nyts' la 19	What sets (CompCert ap	art?							
anatha le aith 1 Llevanai 5 - le' no 189 1 1 - no 19	CompCert is t from miscom C program.	he only produc pilation issues.	tion compiler that The code it produc	is formally verified, tes is proved to beh	using machine-assist ave exactly as specifie	ed mathematical proofs, to d by the semantics of the s	be exempt source			
end.	This level of c levels of softv	onfidence in the are assurance.	e correctness of th	e compilation proce	ss is unprecedented	and contributes to meeting	the highest			
Befinition beanch.map.com forest pc, match cipc with 1 Sowell.breanch 5 :: b) = Unege (firt of) pc =	- 00000						Programmed in Garri PowerPC assembly Programmed and	ariad on	áinymet e	£
pc76+a7 5 - 77 Unage (557 (19) pc + and	The formal pr	pof covers all tr	ansion and a second sec	m the abstract synt	ax tree to the general	ed assembly code. To prec	proved in Coq	dition la	17466- 179 ion in 196 1	Nati 11
serveimage ipeg	~					,			Show al	l ×

Overview About Research Education Industry People Projects Events Publications Institutions Jobs Visitors Student Positions Login

the science of deep specification

Education

DeepSpec is an Expedition in Computing funded by the National Science Foundation.

We focus on the specification and verification of full functional correctness of software and hardware.

Research

We have several major research projects, and our ambitious goal is to connect them at specification interfaces to prove end-to-end correctness of whole systems.

To deliver secure and reliable products, the software industry of the future needs engineers trained in specification and verification. We'll produce that curriculum.







What Are Automated Theorem Provers?

- · Computer programs that (try to) determine if
 - A conjecture C is a logical consequence of a set of axioms Ax
 - · The derivation of conclusions that follow inevitably from facts.
- Systems: Vampire, E, SPASS, Prover9, Z3, CVC4, Satallax, iProver, ...
- Brute-force search calculi (resolution, superposition, tableaux, SMT, ...)
- · Human-designed heuristics for pruning of the search space
- Fast combinatorial explosion on large knowledge bases like Flyspeck and Mizar
- Need to be equipped with good domain-specific inference guidance ...
- ... this what we will try to do here ...
- ... by learning from the knowledge bases and reasoning feedback ...
- Details on particular ATP systems and ML settings later

http://grid01.ciirc.cvut.cz/~mptp/out4.ogv

Using Learning to Guide Theorem Proving

- · high-level: pre-select lemmas from a large library, give them to ATPs
- · high-level: pre-select a good ATP strategy/portfolio for a problem
- high-level: pre-select good hints for a problem, use them to guide ATPs
- low-level: guide every inference step of ATPs (tableau, superposition)
- · Iow-level: guide every kernel step of LCF-style ITPs
- mid-level: guide application of tactics in ITPs
- mid-level: invent suitable ATP strategies for classes of problems
- mid-level: invent suitable conjectures for a problem
- mid-level: invent suitable concepts/models for problems/theories
- · proof sketches: explore stronger/related theories to get proof ideas
- · theory exploration: develop interesting theories by conjecturing/proving
- feedback loops: (dis)prove, learn from it, (dis)prove more, learn more, ...

Sample of Learning Approaches We Have Been Using

- **neural networks** (statistical ML) backpropagation, deep learning, convolutional, recurrent, etc.
- decision trees, random forests, gradient tree boosting find good classifying attributes (and/or their values); more explainable
- **support vector machines** find a good classifying hyperplane, possibly after non-linear transformation of the data (*kernel methods*)
- **k-nearest neighbor** find the *k* nearest neighbors to the query, combine their solutions
- naive Bayes compute probabilities of outcomes assuming complete (naive) independence of characterizing features (just multiplying probabilities)
- inductive logic programming (symbolic ML) generate logical explanation (program) from a set of ground clauses by generalization
- genetic algorithms evolve large population by crossover and mutation
- combinations of statistical and symbolic approaches (probabilistic grammars, semantic features, ...)
- supervised, unsupervised, reinforcement learning (actions, explore/exploit, cumulative reward)

Learning – Features and Data Preprocessing

- Extremely important if irrelevant, there is no use to learn the function from input to output ("garbage in garbage out")
- Feature discovery a big field
- Deep Learning design neural architectures that automatically find important high-level features for a task
- Latent Semantics, dimensionality reduction: use linear algebra (eigenvector decomposition) to discover the most similar features, make approximate equivalence classes from them
- word2vec and related methods: represent words/sentences by *embeddings* (in a high-dimensional real vector space) learned by predicting the next word on a large corpus like Wikipedia
- math and theorem proving: syntactic/semantic patterns/abstractions
- · how do we represent math objects (formulas, proofs, ideas) in our mind?

Neural Autoformalization (Wang et al., 2018)

- · generate about 1M Latex Mizar pairs based on Bancerek's work
- train neural seq-to-seq translation models (Luong NMT)
- evaluate on about 100k examples
- · many architectures tested, some work much better than others
- very important latest invention: attention in the seq-to-seq models
- more data very important for neural training our biggest bottleneck (you can help!)

Rendered LATEX Mizar	If $X \subseteq Y \subseteq Z$, then $X \subseteq Z$.
	X c= Y & Y c= Z implies X c= Z;
Tokenized Mizar	
	X c= Y & Y c= Z implies X c= Z ;
latex	
	If $X \sum Z^{,} \$
Tokenized LATEX	
	If $ X \ y \ Y \ y \ z \ , then X \ y \ z \ .$

Parameter	Final Test	Final Test	Identical	Identical
	Perplexity	BLEU	Statements (%)	No-overlap (%)
128 Units	3.06	41.1	40121 (38.12%)	6458 (13.43%)
256 Units	1.59	64.2	63433 (60.27%)	19685 (40.92%)
512 Units	1.6	67.9	66361 (63.05%)	21506 (44.71%)
1024 Units	1.51	61.6	69179 (65.73%)	22978 (47.77%)
2048 Units	2.02	60	59637 (56.66%)	16284 (33.85%)

Rendered ⊮T _E X	Suppose s_8 is convergent and s_7 is convergent . Then $\lim(s_8+s_7) = \lim s_8 + \lim s_7$
Input LATEX	<pre>Suppose \$ { s _ { 8 } } \$ is convergent and \$ { s _ { 7 } } \$ is convergent . Then \$ \mathop { \rm lim } ({ s _ { 8 } } { + } { s _ { 7 } }) \mathrel { = } \mathop { \rm lim } { s _ { 8 } } { + } \mathop { \rm lim } { s _ { 7 } } \$.</pre>
Correct	<pre>seq1 is convergent & seq2 is convergent implies lim (seq1 + seq2) = (lim seq1) + (lim seq2) ;</pre>
Snapshot- 1000	x in dom f implies (x * y) * (f (x (y (y y)))) = (x (y (y (y y))))) ;
Snapshot- 2000	seq is summable implies seq is summable ;
Snapshot- 3000	seq is convergent & lim seq = 0c implies seq = seq ;
Snapshot- 4000	<pre>seq is convergent & lim seq = lim seq implies seq1 + seq2 is convergent ;</pre>
Snapshot- 5000	<pre>seq1 is convergent & lim seq2 = lim seq2 implies lim_inf seq1 = lim_inf seq2 ;</pre>
Snapshot- 6000	<pre>seq is convergent & lim seq = lim seq implies seq1 + seq2 is convergent ;</pre>
Snapshot- 7000	seq is convergent & seq9 is convergent implies lim (seq + seq9) = (lim seq) + (lim seq9) ;